

Growth Rate of the Binomial Coefficient

Hanan Ayad

1 Computing $C(n, k)$

The number of additions made by the recursive algorithm for computing $C(n, k)$ based on the formula $C(n, k) = C(n - 1, k - 1) + C(n - 1, k)$ is equal to $C(n, k) - 1$, which can be said to be $\Theta(C(n, k))$. But, what is the growth rate of $C(n, k)$? Specifically, how is $C(n, k)$ expressed in terms of the basic asymptotic efficiency classes?

2 Growth Rate of $C(n, k)$

To express the growth rate of $C(n, k)$ in terms of n , we consider the function $C(n, k)$ for different values of k .

For $k = 1$,

$$C(n, 1) = n \in \Theta(n).$$

For $n = 2$,

$$C(n, 2) = \frac{n(n-1)}{2} \in \Theta(n^2).$$

For $n = 3$,

$$C(n, 3) = \frac{n(n-1)(n-2)}{6} \in \Theta(n^3).$$

For $n = \frac{n}{2}$,

$$C(n, n/2) = \frac{n!}{((n/2)!)^2}.$$

Using Stirling formula: $n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$, as $n \rightarrow \infty$, we get,

$$C(n, n/2) = \frac{\sqrt{2\pi n} (n/e)^n}{[\sqrt{2\pi n/2} (n/2e)^{n/2}]^2} = \frac{2}{\sqrt{2\pi n}} 2^n \in \Theta(2^n n^{-0.5}) \quad (1)$$

Equation 1 implies that $C(n, k)$ grows very fast for k close to $n/2$